

# DSC 190

DATA STRUCTURES & ALGORITHMS

Lecture 15 | Part 1

**Today's Lecture**

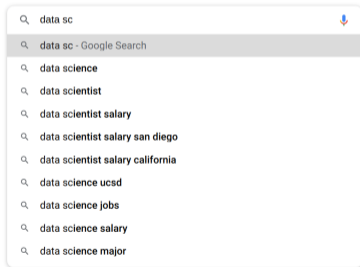
# String Data Structures

- ▶ One of the themes of this quarter:
- ▶ If you're doing something once, use an algorithm.
- ▶ If you're doing it over and over, use an appropriate data structure.

# String Data Structures

- ▶ Over the next two lectures, we'll look at data structures for strings.
- ▶ Today: **tries** for efficient repeated **prefix queries**.

# Autocompletion



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Lecture 15 | Part 2

**Tries**

# Trie

- ▶ A data structure for storing strings.
- ▶ Pronounced “try”, short for “re**trieval**”.
- ▶ Supports fast **prefix query** and **membership query**.

# Prefixes

- ▶ A **prefix**  $p$  of a string  $s$  is a contiguous slice of the form  $s[0:t]$ , for some  $t$ .
- ▶ Examples:
  - ▶ "test" is a prefix of "testing"
  - ▶ "te" is a prefix of "testing"
  - ▶ "sa" is a prefix of "san diego"
  - ▶ "di" is **not** a prefix of "san diego"

# Prefix Query

- ▶ **Given:** a collection of  $n$  strings and a prefix,  $p$ .
- ▶ **Find:** all strings in the collection for which  $p$  is a prefix.
- ▶ Example:
  - ▶ "bar", "bad", "bid", "car"
  - ▶  $p = \text{"ba"}$

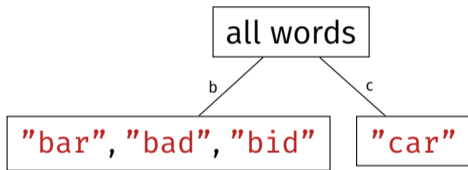


# Brute Force

- ▶ Loop over each of  $n$  strings, compare against prefix  $p$ .
- ▶ Worst-case time:  $\Theta(n \cdot |p|)$

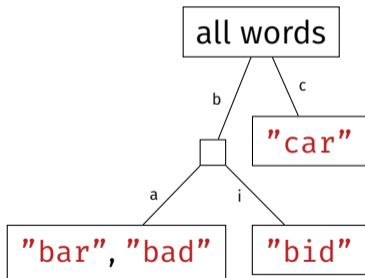
# Trie: Motivation

"bar", "bad", "bid", "car"



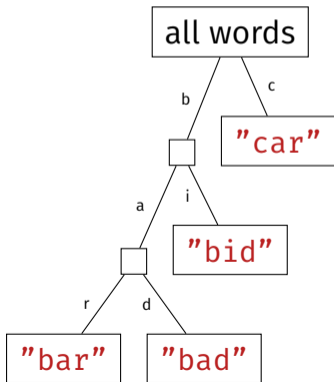
# Trie: Motivation

"bar", "bad", "bid", "car"

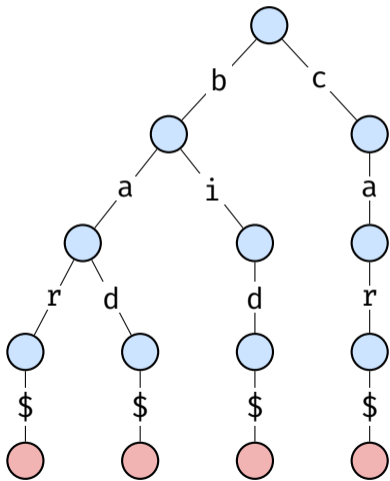


# Trie: Motivation

"bar", "bad", "bid", "car"

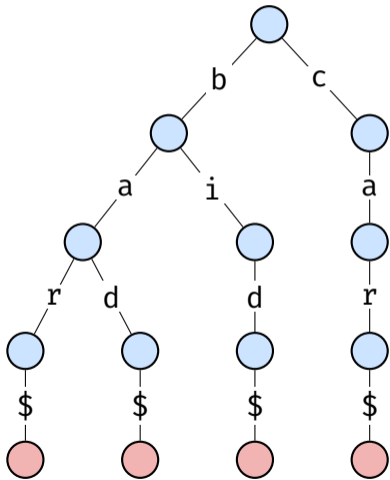


# Tries



- ▶ Internal nodes represent prefixes.
- ▶ Leaf nodes represent full words.
- ▶ Edges are characters.
- ▶ Words are encoded as paths.

# Sentinels



- ▶ \$ is a **sentinel**.
- ▶ It is different from the dollar sign character.
- ▶ It marks the end of a word.
- ▶ Used to show that "bar" in trie, but "ba" not.

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**Implementing Tries**

# Representation

- ▶ Each node has a hash table / array mapping characters to a child nodes.
- ▶ Sentinel represented with a singleton object?

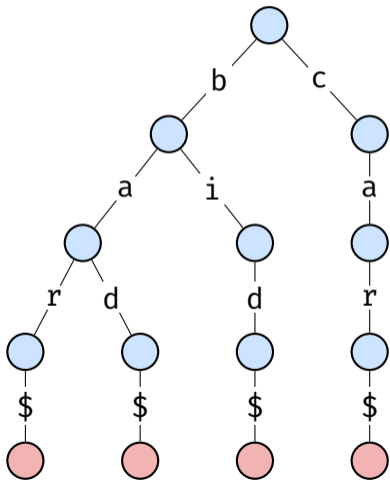
```
END_OF_STRING = object()
```

```
class TrieNode:
```

```
    def __init__(self):  
        self.children = {}
```

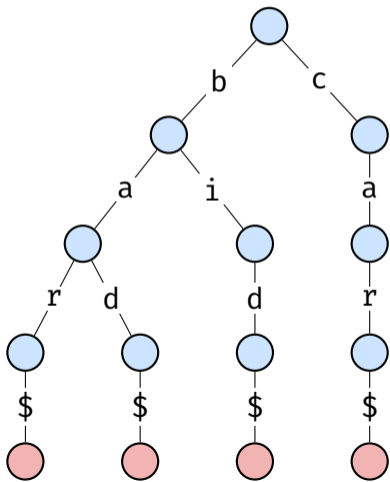


# Insertion



- ▶ “Walk” down tree, creating edges and nodes as necessary.
- ▶ When no more letters left, add sentinel.
- ▶ Example: insert  
“cab”, “card”, “zoo”

# Insertion (Recursive)



- ▶ Suppose we `.insert(s)` on root node.
- ▶ If `s[0]` not in `self.children`, create a new node.
- ▶ Otherwise, let child be `self.children[s[0]]`.
- ▶ Recursively insert `s[1:]` into child.

```
def insert(self, s, start=0, stop=None):
    """Insert s[start:stop] into the trie."""
    if stop is None:
        stop = len(s)

    if start >= stop:
        self.children[END_OF_STRING] = TrieNode()
        return

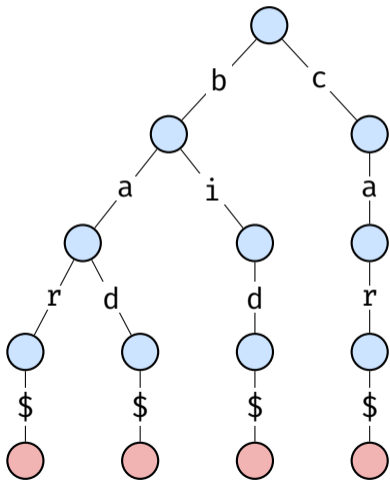
    if s[start] not in self.children:
        self.children[s[start]] = TrieNode()

    child = self.children[s[start]]
    child.insert(s, start + 1, stop)
```

# Insertion Time Complexity

- ▶  $\Theta(|w|)$  time, where  $w$  is the string inserted.
- ▶ No matter how many elements in trie!

# Walk



- ▶ Useful operation.
- ▶ Given a prefix, “walk” down tree.
- ▶ If we “fall off”, raise error.
- ▶ Otherwise, return last node seen.
- ▶ Examples: **“ba”**, **“bo”**

```
def walk(self, s, start=0, stop=None):
    """Walk the trie following s[start:stop].
    Raises ValueError if falls off tree.
    Returns last node encountered otherwise."""
    if stop is None:
        stop = len(s)

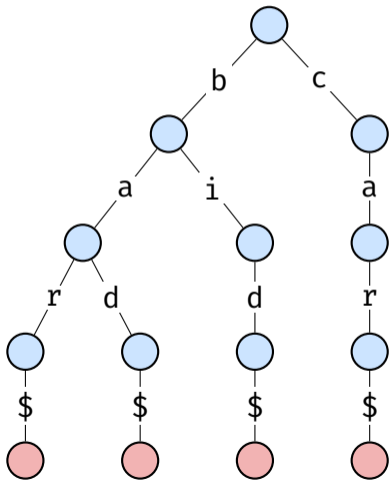
    if start >= stop:
        return self

    if s[start] not in self.children:
        raise ValueError('Fell off tree.')
    else:
        child = self.children[s[start]]
        return child.walk(s, start + 1, stop)
```

# Walk Time Complexity

- ▶ Worst-case  $\Theta(|p|)$  time, where  $p$  is the prefix searched.
- ▶ No matter how many elements in trie!

# Membership Query



- ▶ Given  $p$ , return True/False if  $p$  in collection.
- ▶ “Walk” down tree.
- ▶ If we “fall off”, return False.
- ▶ If not, check that sentinel in children.
- ▶ Examples: “ba”, “bad”



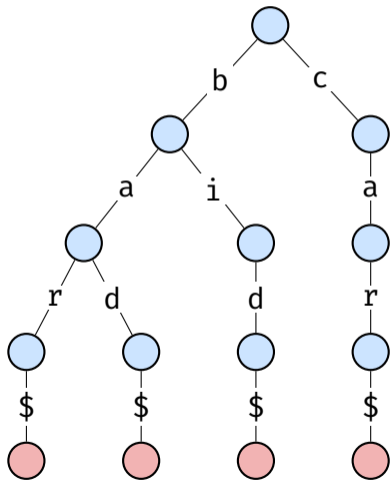
```
def membership_query(self, s, start=0, stop=None):
    """Determine if s[start:stop] is in trie."""
    try:
        node = self.walk(s, start, stop)
    except ValueError:
        return False

    return END_OF_STRING in node.children
```

# Membership Query Time Complexity

- ▶ Worst-case  $\Theta(|w|)$  time, where  $w$  is the prefix searched.
- ▶ No matter how many elements in trie!

# Produce



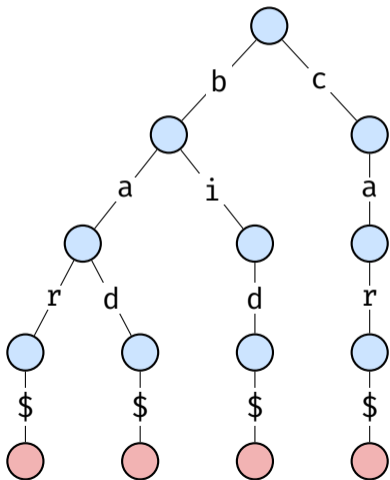
- ▶ Goal: generate all words in subtree.
- ▶ Perform a DFS, keeping track of letters along path.
- ▶ If we find a sentinel, print path.

```
def produce(self, pathchars=''):
    """Generate the words in the trie."""
    for letter, child in self.children.items():
        if letter is END_OF_STRING:
            yield pathchars
        else:
            yield from child.produce(pathchars + letter)
```

## Produce Time Complexity

- ▶ Worst-case  $\Theta(\ell)$  time, where  $\ell$  is total length of all strings stored in the trie.
- ▶ If length strings is considered a constant, this is  $\Theta(n)$ .

# Prefix Query (Complete)



- ▶ Given p, return all completions.
- ▶ “Walk” down tree.
- ▶ If we “fall off”, return empty list..
- ▶ If not, produce all nodes in subtree.
- ▶ Examples: **“ba”**, **“bad”**

```
def complete(self, prefix):
    try:
        node = self.walk(prefix)
    except ValueError:
        return []
    return list(node.produce())
```

## Prefix Query Time Complexity

- ▶ Worst-case  $\Theta(|p| + \ell_p)$  time, where  $p$  is the prefix searched and  $\ell_p$  is the total length of all matches.
- ▶ If length is considered constant, this is  $\Theta(|p| + z)$ , where  $z$  is number of matches.



# DSC 190

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**Demo**